

How do Performance Indicator Parametrizations Influence the Assessment of Algorithm Portfolios?

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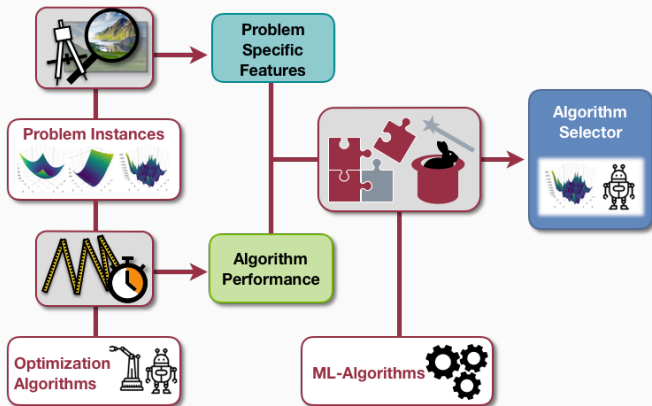
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Introduction

Algorithm Selection

Idea of Algorithm Selection:

- *Algorithm Selection Problem*¹: find the individually best suited algorithm for an unseen optimization problem



1. Rice, J. (1976). *The Algorithm Selection Problem*. In *Advances in Computers* (pp. 65-118).

Requirements:

- Comprehensive benchmark of portfolio solvers (as a foundation for algorithm selection)
- Suitable performance measure needed, e.g., PAR10², ERT³.
- Performance measures often parameterized.
~> How do parameters affect the benchmark results?

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2. Bischl, B. et al. (2016). *ASlib: A Benchmark Library for Algorithm Selection*. In *Artificial Intelligence Journal* (pp. 41-58).
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Our contribution:

- Systematic analysis of parameterizations on a comprehensive benchmark study of inexact TSP solvers.

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- Time limit / cutoff time $T \in \mathbb{R}_{>0}$.

Performance Measures

Numeric Example:

- 10 runs of solvers X and Y
- budget for runtime r_s of successful runs is set to $T = 1$
- solver X: 8 successful runs ($p_f = 0.2$) with $r_s = 0.8$
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- ~> How do we aggregate the runs (meaningfully)?

Penalized Average Runtime (PAR)⁴:

- Arithmetic mean of running times, $r_i^{A,l}, i \in [m]$
- Expired runs are penalized by $f \cdot T$, where f is the **penalty factor**

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$$\text{PAR}_{A,l}(f) := \frac{1}{m} \sum_{i=1}^m \tilde{r}_i^{A,l} \quad \text{with} \quad \tilde{r}_i^{A,l} = \begin{cases} f \cdot T, & \text{if } r_i^{A,l} > T \\ r_i^{A,l}, & \text{otherwise.} \end{cases}$$

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- Usually, the rather heuristic value $f = 10$ is employed.
 $\leadsto \text{PAR}_{A,l}(10)$

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Penalized Quantile Runtime (PQR)⁵:

- Replace outlier-sensitive mean by more robust p -quantile, $p \in (0, 1]$.

5. Kerschke, P. et al. (2018). *Parameterization of State-of-the-Art Performance Indicators: A Robustness Study Based on Inexact TSP Solvers*. In Proceedings of GECCO 2018 Companion (pp. 1737-1744).

Performance Measures (cont.)

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$$\text{PQR}_{A,l}(p, f) := \begin{cases} f \cdot T, & \text{if } \sum_{i=1}^m \mathbb{1}\{r_i^{A,l} < T\} < \lfloor mp + 1 \rfloor \\ q_p(r_1^{A,l}, \dots, r_m^{A,l}), & \text{otherwise.} \end{cases}$$

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Expected Runtime (ERT)⁶:

- Popular / most common measure in continuous optimization.

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Performance Measures (cont.)

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$$\begin{aligned} \text{ERT}_{A,l} &= \frac{1}{s} \sum_{j=1}^s r_{i_j}^{A,l} + \left(\frac{1 - \rho_s}{\rho_s} \right) \cdot T \\ &= \frac{1}{s} \left(\sum_{j=1}^s r_{i_j}^{A,l} + (m - s) \cdot T \right). \end{aligned}$$

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- Corresponds to average runtime for observing one successful run.

6. Hansen, N. et al. (2009). *Real-Parameter Black-Box Optimization Benchmarking 2009: Experimental Setup*. In INRIA Research Report RR-6828.

Penalized Expected Runtime (PERT):

- Introducing penalty factor into ERT.

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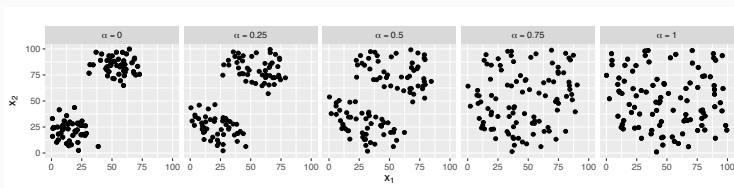
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Performance Indicator	$f = 10$		$f = 100$	
	X	Y	X	Y
$PAR_{A,I}(f)$	2.64	8.04	20.64	80.04
$PQR_{A,I}(0.5, f)$	0.80	10.00	0.80	100.00
$ERT_{A,I}$	1.05	4.20	1.05	4.20
$PERT_{A,I}(f)$	3.30	40.20	25.80	400.20

Case Study

Case Study

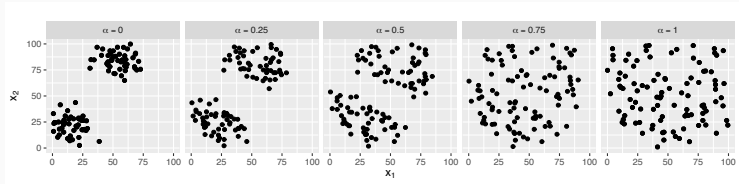
- Based on performance data from our previous TSP algorithm selection study⁷:
- Five state-of-the-art inexact TSP solvers (**Algorithms \mathcal{A}**):
 - MAOS [4], EAX [3], LKH [2], EAX+restart and LKH+restart [1].
- Six sets of TSP instances (**Problems \mathcal{I}**):
 - VLSI, TSPLIB, National, RUE, clustered (netgen) and morphed.



7. Kerschke, P. et al. (2017). *Leveraging TSP Solver Complementarity through Machine Learning*. In ECJ.

Case Study

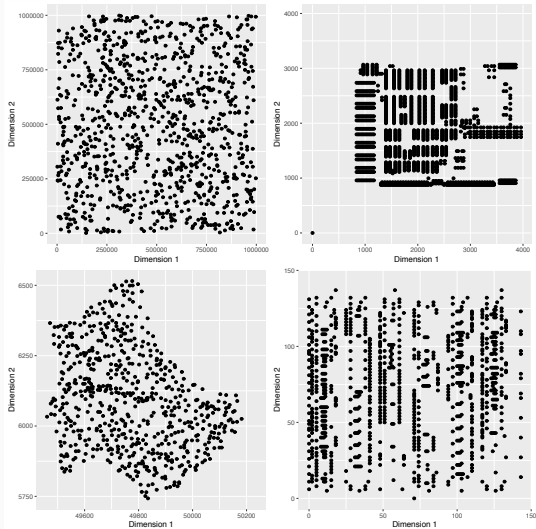
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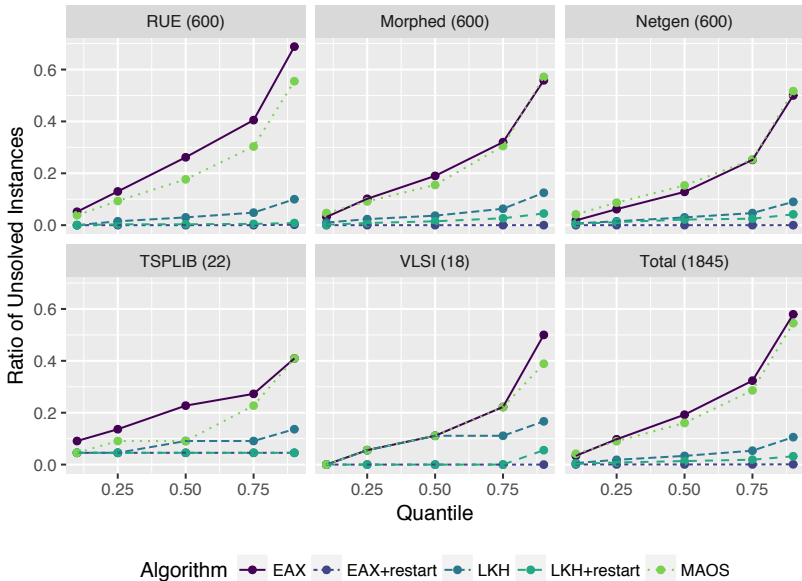
- EAX+restart was single-best-solver (SBS) regarding $PAR_{A,I}(10)$.

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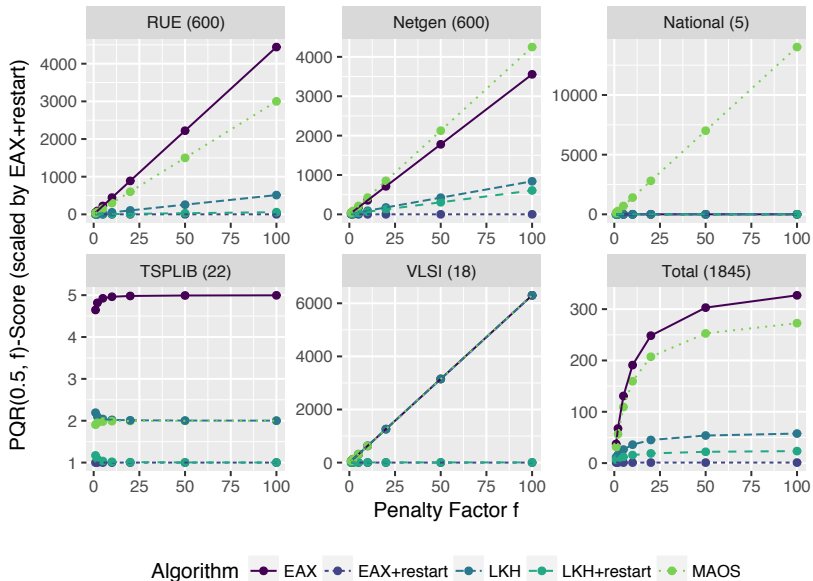
Case Study



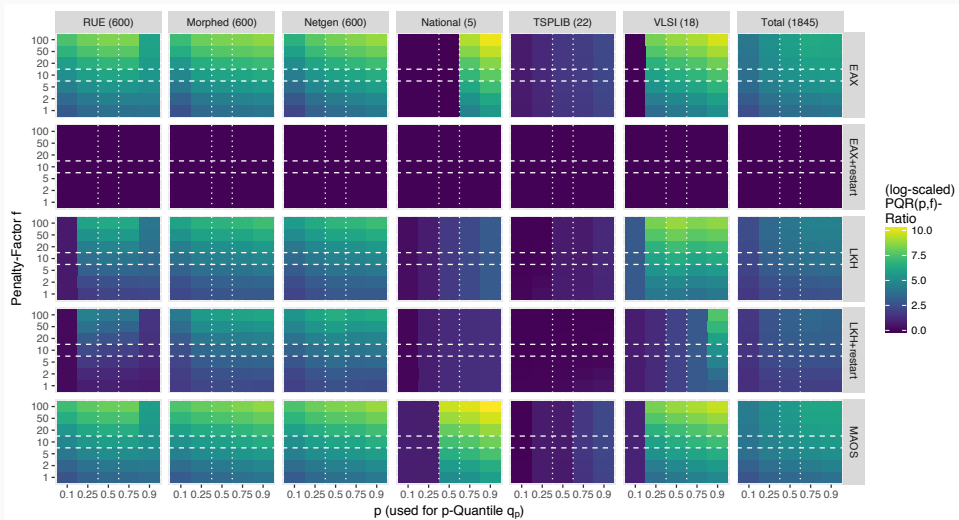
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Observations:

- Finding a suitable pair of penalty factor f and quantile p quickly becomes very complex.
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Idea:

- One basically wants to optimize the runtime and success probability simultaneously.
- Why not use a multi-objective approach?
 \leadsto consider for instance HV principle

Case Study

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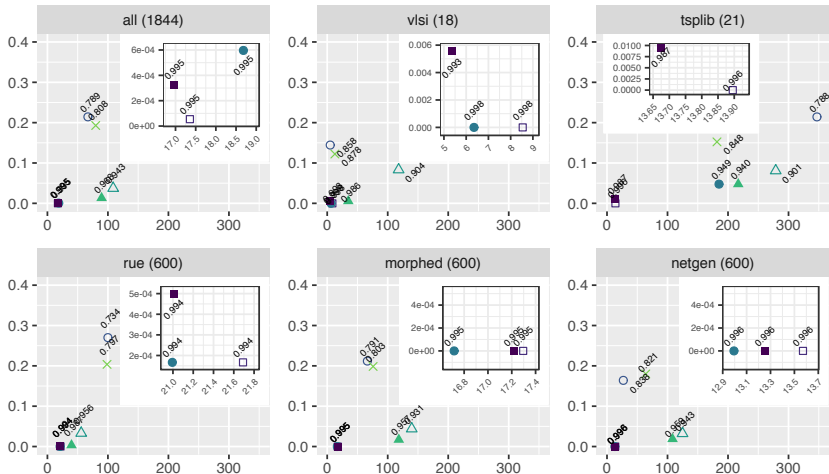
$$HV_{A,l} = (T - r_s) \cdot (1 - p_f).$$

Performance Indicator	$f = 10$		$f = 100$	
	X	Y	X	Y
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$HV_{A,l}$	0.16	0.16	0.16	0.16

Case Study

(Visual and Measure Independent) Comparison of TSP Solvers:

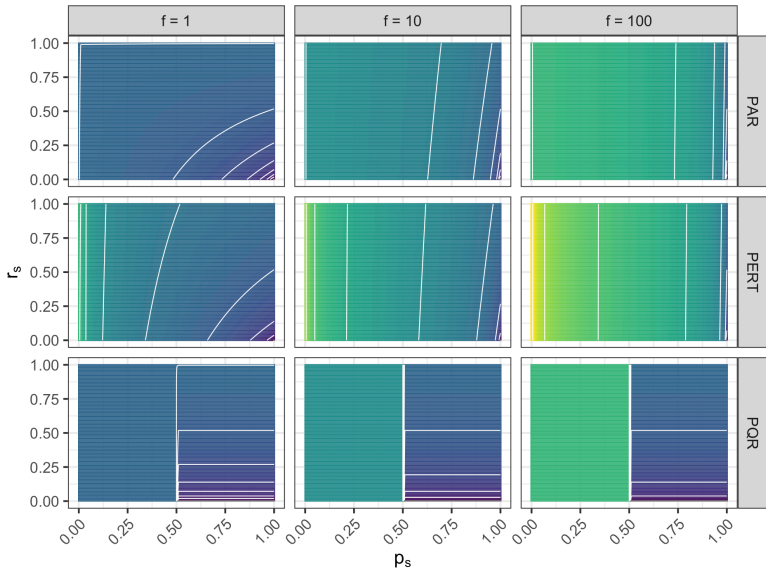
Algorithm ■ AS-ECJ □ AS-UBC ○ EAX ● EAX+restart ▲ LKH ▲ LKH+restart × MAOS



Further Comparison of Performance Measure Parameterizations

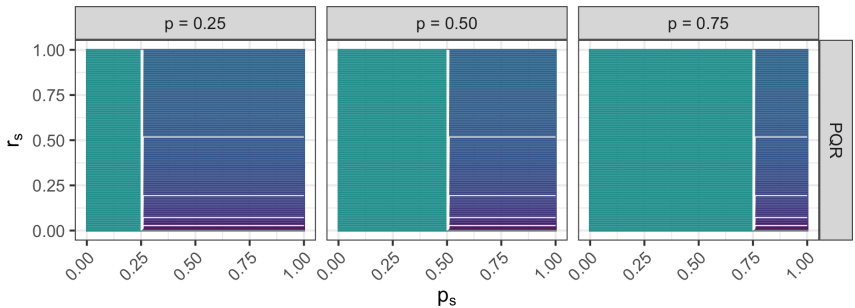
Further Comparison of Performance Measure Parameterizations

(Theoretical) Effect of Penalty Factor on Performance Measures:



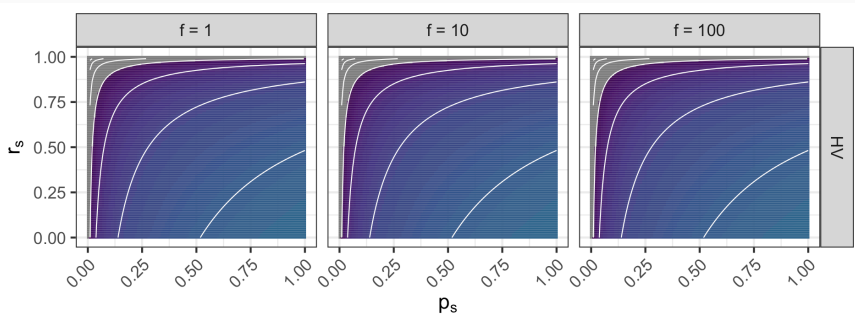
Further Comparison of Performance Measure Parameterizations

(Theoretical) Effect of Quantile on PQR(p , 10)-Score:



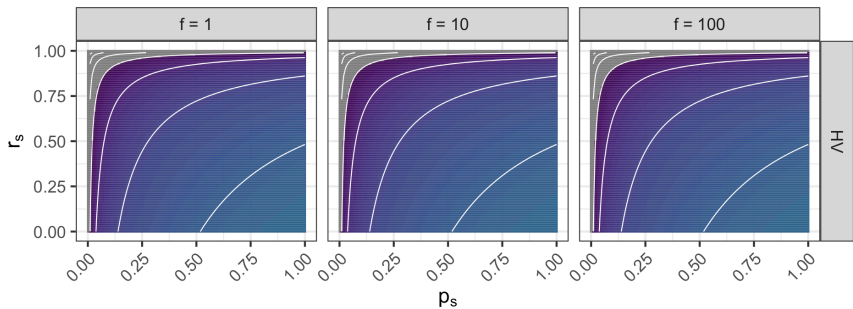
Further Comparison of Performance Measure Parameterizations

HV indicator as performance measure:



Further Comparison of Performance Measure Parameterizations

HV indicator as performance measure:



Note that HV is robust against alterations of the penalty score.

Conclusion & Outlook

Conclusion:

- We systematically analyzed effects of different parameterizations of performance indicators.
- Varying penalty factor allows for altering leverage of failed runs.
- (P)ERT is much more prone to single runs
 \leadsto huge impact of single failed runs.
- Choosing a suitable measure has a huge impact on the actual performance assessment (for solvers and selectors).
- HV might be a good alternative to common measures.

Outlook:

- Theoretical investigations of indicators.
- Introduction of alternative (multi-objective) indicators⁸.
- Application in context of algorithm selection.

8. Bossek, J. & Trautmann, H. (2018). *Multi-Objective Performance Measurement: Alternatives to PAR10 and Expected Running Time*. In Proceedings of LION 2018.

Questions?

References

- [1] Jérémie Dubois-Lacoste, Holger H. Hoos, and Thomas Stützle. On the Empirical Scaling Behaviour of State-of-the-art Local Search Algorithms for the Euclidean TSP. In **Proceedings of the 17th Annual Conference on Genetic and Evolutionary Computation (GECCO)**, pages 377 – 384. ACM, 2015. ISBN 978-1-4503-3472-3. doi: 10.1145/2739480.2754747. URL <http://dl.acm.org/citation.cfm?id=2754747>.
- [2] Keld Helsgaun. General k-opt Submoves for the Lin-Kernighan TSP Heuristic. **Mathematical Programming Computation**, 1(2-3):119 – 163, 2009. doi: 10.1007/s12532-009-0004-6. URL <https://link.springer.com/article/10.1007/s12532-009-0004-6>.
- [3] Yuichi Nagata and Shigenobu Kobayashi. A powerful genetic algorithm using edge assembly crossover for the traveling salesman problem. **INFORMS Journal on Computing**, 25(2):346 – 363, 2013. doi: 10.1287/ijoc.1120.0506. URL <http://dl.acm.org/citation.cfm?id=2466704>.
- [4] Xiao-Feng Xie and Jiming Liu. Multiagent Optimization System for Solving the Traveling Salesman Problem (TSP). **IEEE Transactions on Systems, Man, and Cybernetics, Part B: Cybernetics**, 39(2):489 – 502, 2009. URL <http://ieeexplore.ieee.org/document/4717264/>.